

# Time-varying resilient virtual network mapping for multi-location cloud data centers

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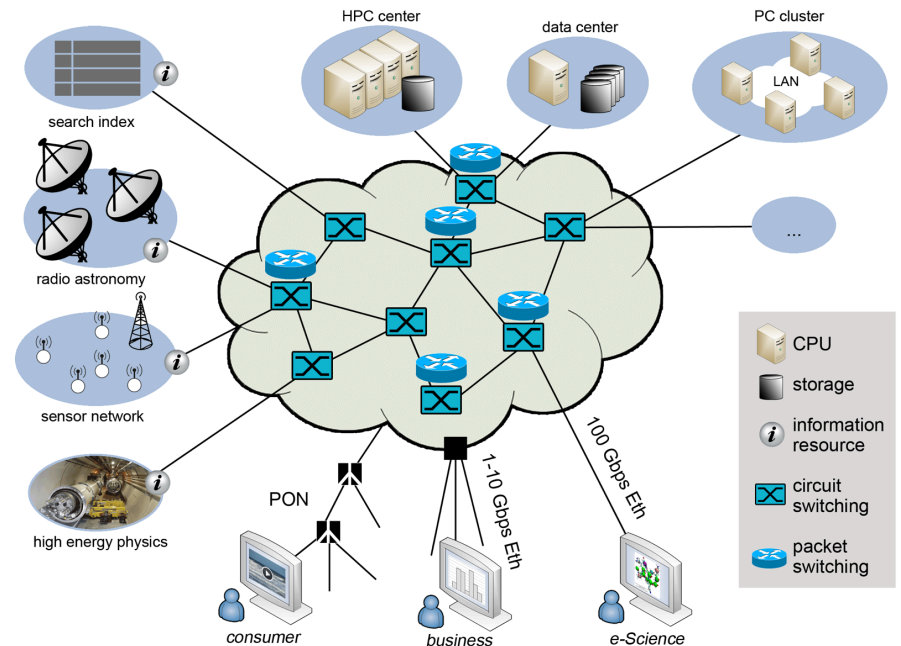
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3: INTEC – IBCN, Ghent University – iMinds, Ghent, Belgium

# Optical clouds

Optical networks crucial for increasingly demanding cloud services, e.g.,

- Computing:
  - High energy physics
  - Amazon EC2, Microsoft Azure
- Online storage:
  - Dropbox, Google Drive, etc.
- Collaboration tools:
  - MSOffice 365, Google Docs
- Video streaming:
  - Netflix, YouTube

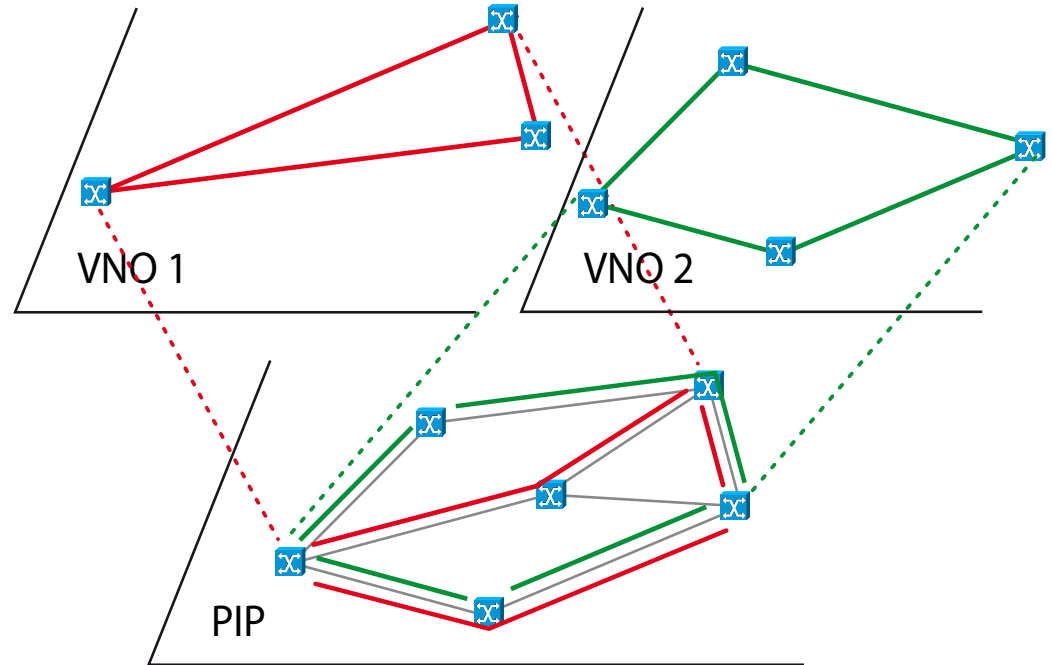


C. Develder, et al., "Optical networks for grid and cloud computing applications", Proc. IEEE, Vol. 100, No. 5, May 2012, pp. 1149-1167.

# Network virtualization

Physical network is logically partitioned in isolated virtual networks

- **Virtual Network Operators (VNO)** operate logically separate networks
- **Physical Infrastructure Providers (PIP)** have full control over infrastructure (fibers, OXCs)



J.A. García-Espín, et al., "Logical Infrastructure Composition Layer: the GEYSERS holistic approach for infrastructure virtualisation", in Proc. TERENA Networking Conference (TNC 2012), Reykjavík, Iceland, 21-24 May 2012.

# Overview

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2. Problem statement
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# Resiliently provisioning virtual cloud networks

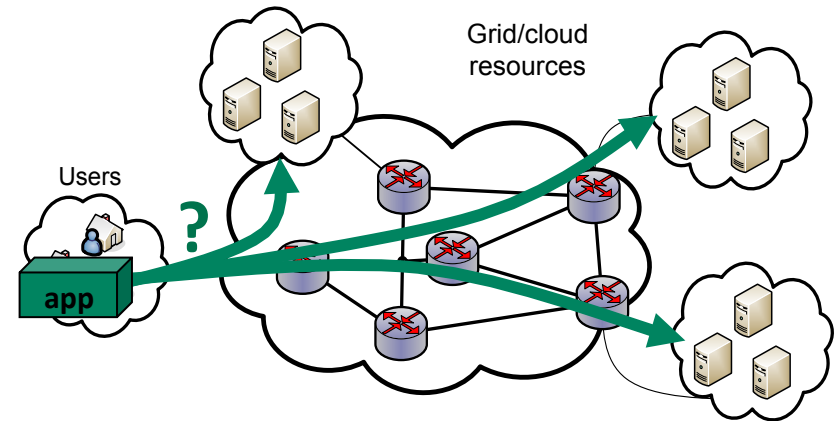
How to choose the virtual to physical mapping, such that

- Services remain available in case of network failures
- Bandwidth for providing services is minimal

?

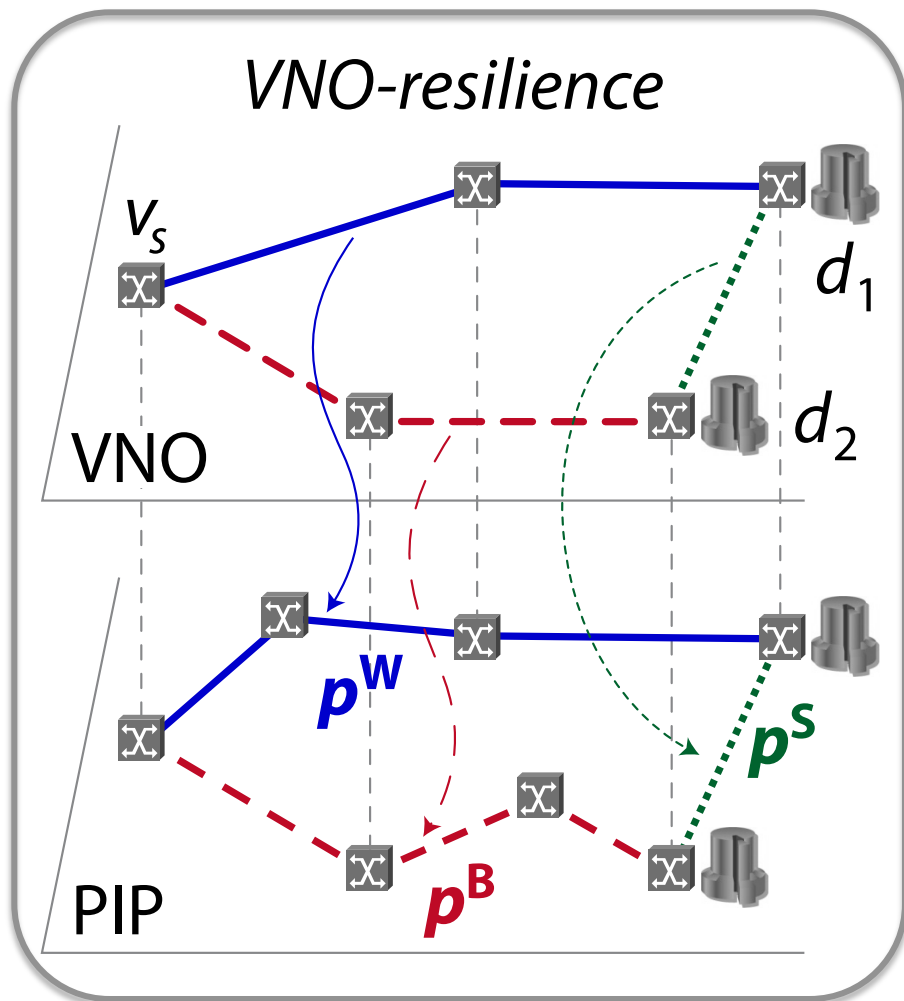
Note:

- **Anycast**: requests coming from users can be served by any server
- Cloud services offered by VNO
- Cloud services run on top of PIP

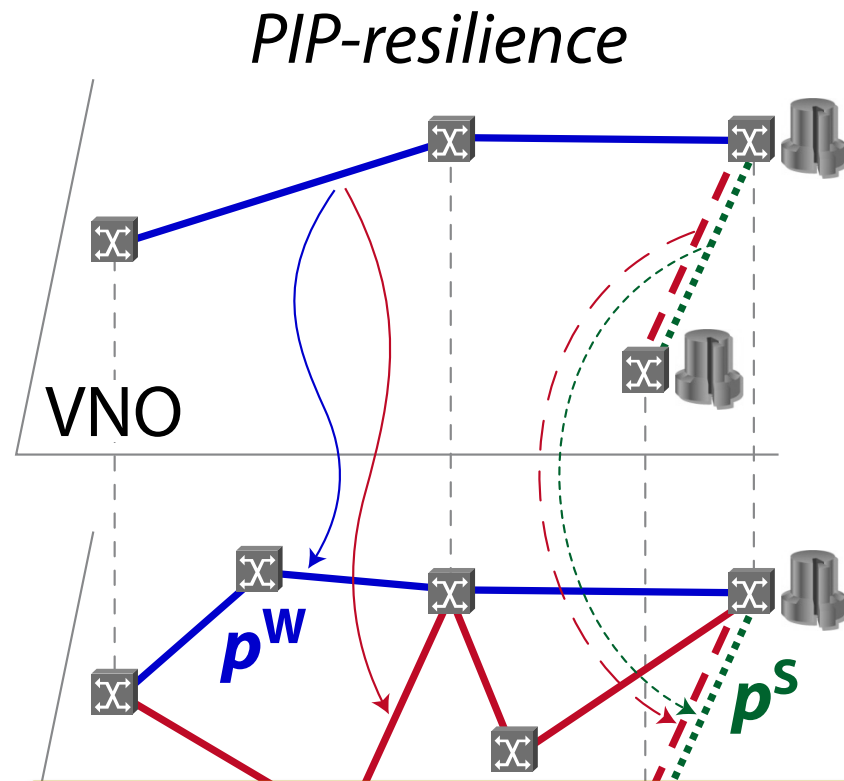


B. Jaumard, A. Shaikh and C. Develder, "Selecting the best locations for data centers in resilient optical grid/cloud dimensioning (Invited Paper)", in Proc. 14th Int. Conf. Transparent Optical Netw. (ICTON 2012), Coventry, UK, 2-5 Jul. 2012.

# Two proposed protection schemes:



*This paper*



M. Bui, B. Jaumard, and C. Develder, "Anycast end-to-end resilience for cloud services over virtual optical networks" (Invited Paper), in Proc. 15th Int. Conf. Transparent Optical Netw. (ICTON 2013), Cartagena, Spain, 23-27 Jun. 2013.

# Related work: Static traffic scenarios

- **Traditional dimensioning (no virtualisation, no resilience):**
  - Develder *et al.* 2009: Anycast, flexibility in choosing data center
- **Resilient dimensioning problem:**
  - Shaikh *et al.* 2011, Develder *et al.* 2013: scalable method, no synchronization between working and backup DCs
- Routing cloud service requests and **mapping a VNet** to the physical infrastructure separately:
  - Lee *et al.* 2009, Yu *et al.* 2010: Survivable VNet embedding, but *assume VNet is given*
  - Jiang *et al.* 2012, Alicherry *et al.* 2012: Optimal server selection and routing of anycast services in the physical layer for intra- and inter-DC networks but *no resilient network design in the virtual layer*
- **VNet planning problem:**
  - Barla *et al.* 2012, Barla *et al.* 2013: using mixed integer linear programming, but *no synchronization between working and backup DCs*
  - Bui *et al.* 2013 (ICTON): first model that incorporates synchronisation path, but *still static traffic!*



# Problem statement

- Study **time-varying traffic**:  
Traffic pattern changes from one period ( $t$ ) to the next ( $t+1$ )
- Key research question:  
Benefit (in network resource usage) of **changing routes** for legacy traffic, i.e., that continues from  $t$  to  $t+1$ ?
  - Does it help to only change backup paths?
  - Or do we need to change working as well?
  - For *all* legacy traffic?

# Problem statement

Given:

- Cloud network topology:  $G = (V, L)$ , with  $V =$  nodes,  $L =$  links
- Locations of the data centers,  $V_D \subseteq V$
- Set of service requests,  $K$  partitioned into
  - $K^{LEG}$ : requests in period  $t$ , that continue into the next  $t+1$
  - $K^{ADD}$ : requests in subsequent period  $t+1$and characterized by
  - $v_k$ : source of service
  - $\Delta_k$ : bandwidth requirement in period  $t$
  - Services originating from the same source are aggregated
- Routing of the requests in period  $t$ :  $K^{LEG}$

Find:

- Choice of primary and backup **DC locations** for each service in period  $t+1$
- Primary, backup and synchronization **paths** in period  $t+1$

Such that:

- Total network **bandwidth utilization is minimized**
- $K^{LEG}$  are (i) unchanged, (ii) only changed for backup, (iii) freely changed

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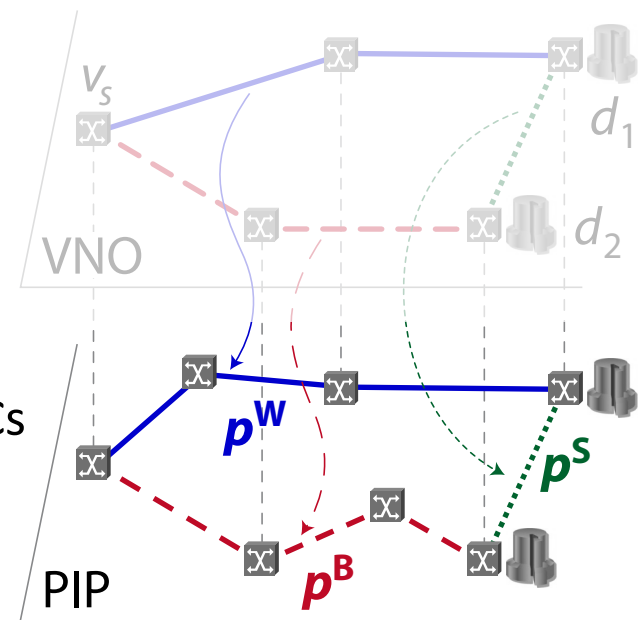
# Solution: Column generation model

- Column generation idea:

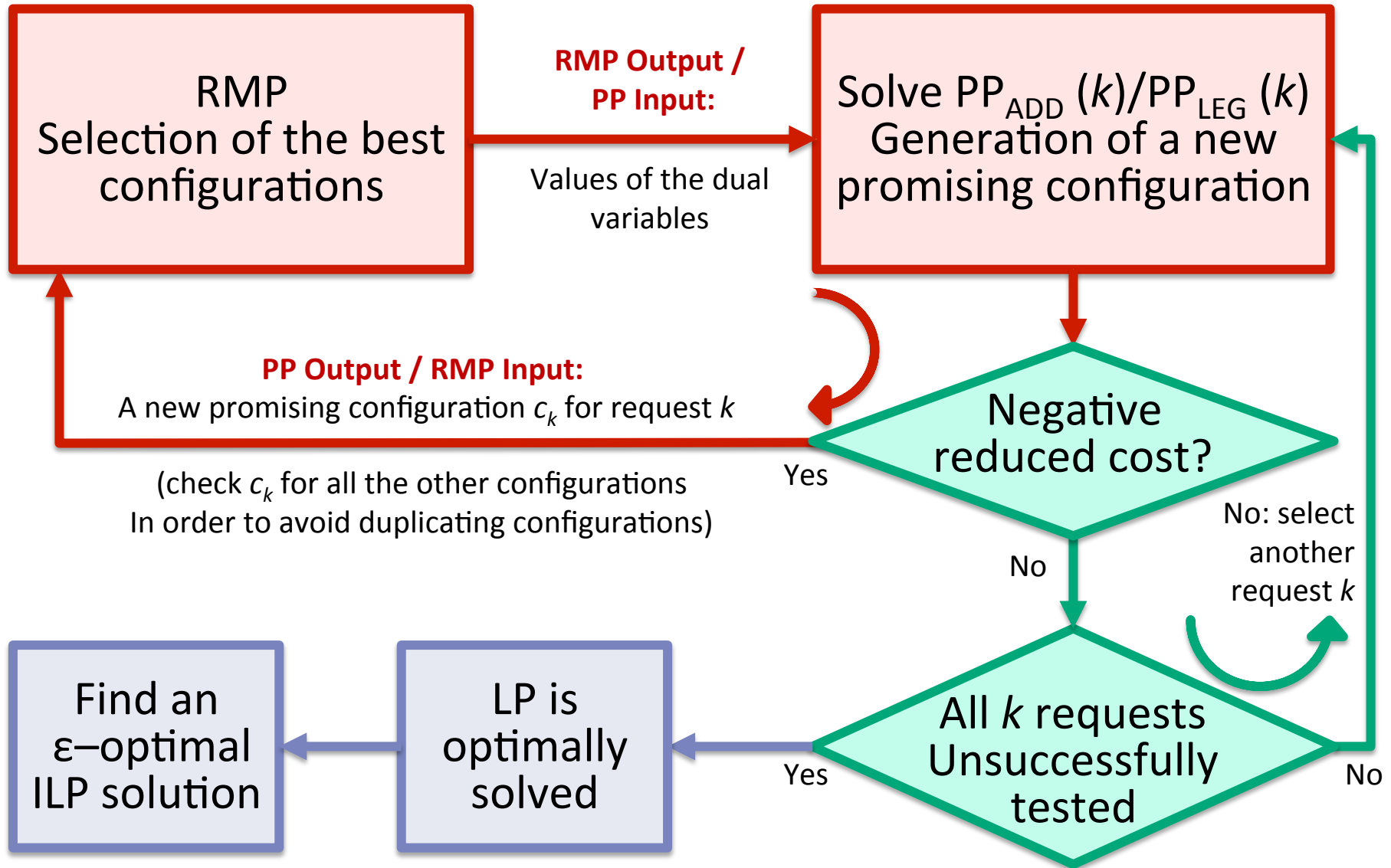
- Many different “configurations”
- Start from a restricted subset of such “configurations”
- Iteratively find additional configurations that reduce the cost:
  - (1) Restricted Master Problem (RMP)
  - (2) Pricing Problem (PP) to find new configs

- A configuration =

- **Working** path
- **Backup** path
- **Sync** path between the primary & backup DCs



# Column generation solution algorithm

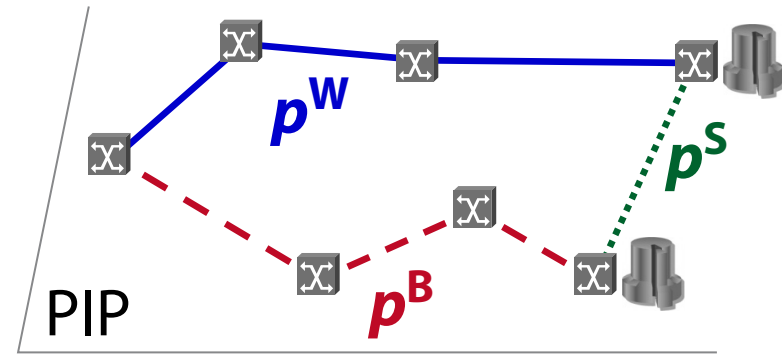


# Restricted Master Problem (RMP)

$$\min \sum_{l \in L} (\underbrace{\beta_l^W + \beta_l^B + \beta_l^S}_{BW_l}) \cdot \|l\|$$

$$+ \text{PENAL}^{\text{DISRUPT\_B}} \sum_{k \in K^{\text{LEG}}} x_k^{\text{LEG\_BS}}$$

$$+ \text{PENAL}^{\text{DISRUPT\_W}} \sum_{k \in K^{\text{LEG}}} x_k^{\text{LEG\_W}}$$



Case (ii):  
penalty for disrupting B/S  
path of legacy

Case (iii):  
penalty for disrupting W  
path of legacy

## Constraints:

- Assure all requests are granted
- Count legacy changes  $x^{\text{LEG\_BS}}$ ,  $x^{\text{LEG\_W}}$
- Compute W, B, S bandwidths
- Check capacity constraints on data centers

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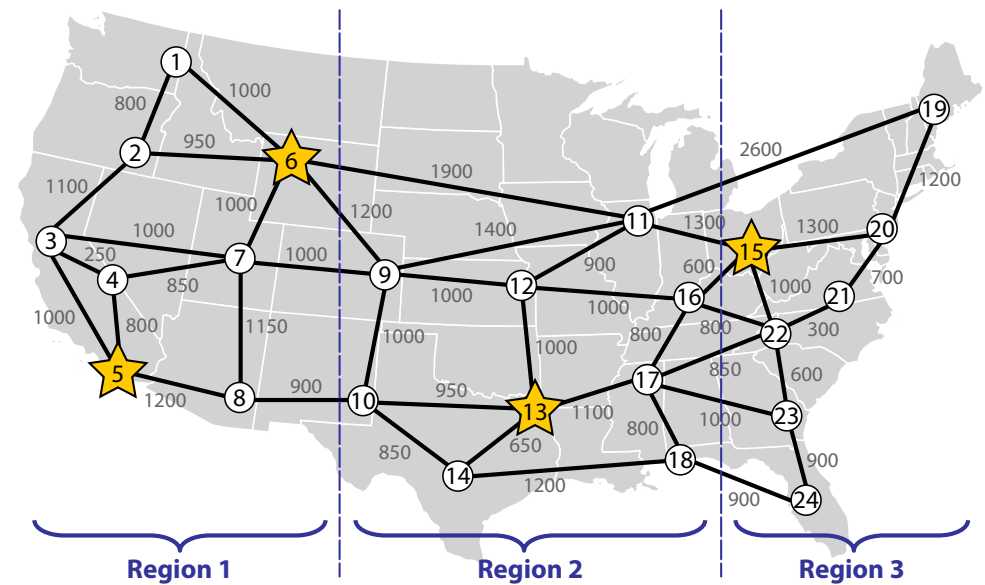
# Case study

## ■ Topology:

- 24 nodes, 43 links
- Data centers in ★: CA, WY, TX, OH

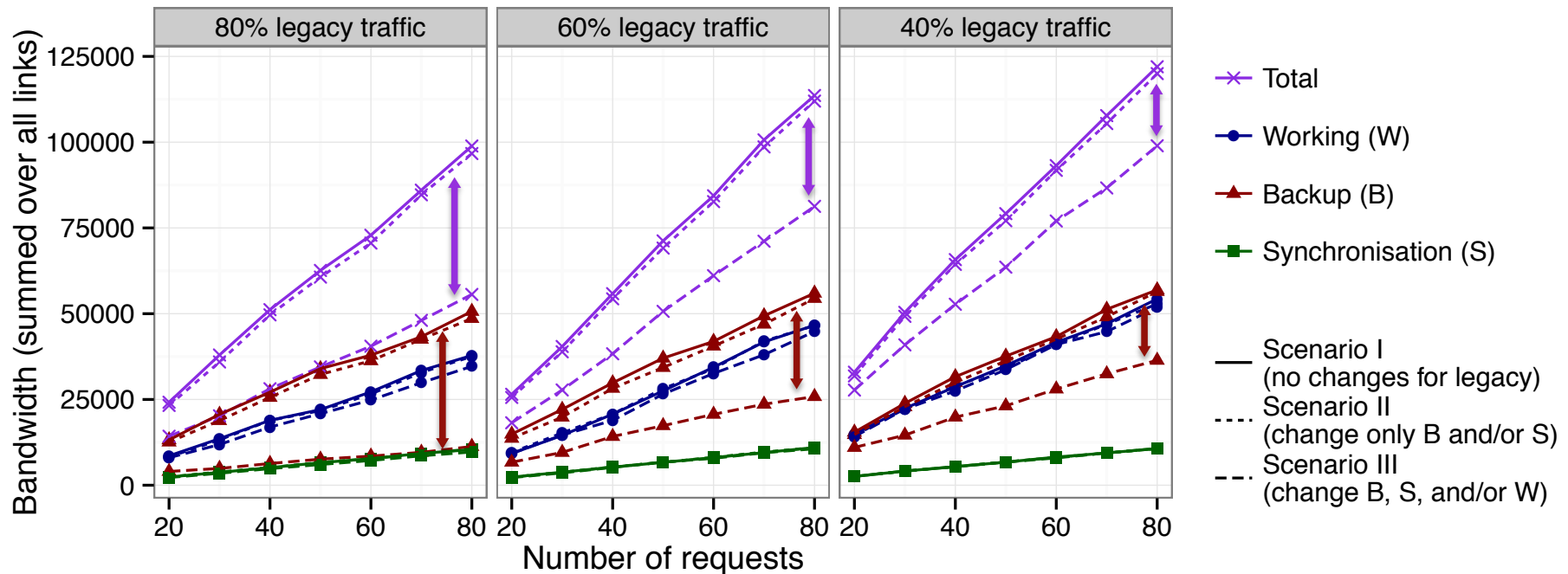
## ■ Traffic:

- **Total** of 20...80 requests (same for  $t$  and  $t+1$ ) with  $\Delta_k$  in  $[0,1]$  units; synchronisation fraction  $\delta_k = 0.1$
- **Period  $t$** : 30% region 1, 50% region 2, 20% region 3
- **Period  $t+1$** :



Scenario	Region 1	Region 2	Region 3
40% legacy	20% drop 20% add	30% drop 10% add	10% drop 30% add
60% legacy	10% drop 10% add	20% drop 10% add	10% drop 20% add
80% legacy	10% drop 10% add	10% drop -	- 10% add

# Results: Net total bandwidth savings?



1. Only changing backup (II, ----) does not **save** much; changing also working (III, - - -) does
2. Capacity savings are realized through **sharing of backup**
3. Savings obviously diminish with decreasing legacy fraction (L→R)

# Results: Change all legacy routes?

Of the legacy requests, we typically change >50%, but only about **20%** need to **reroute the working path**

Legacy traffic	Path changes	Total demand (number of requests)							$\alpha_{LEG}$ over all demands
		20	30	40	50	60	70	80	
80%	Scen. II – Changed B and/or S	9.6	12.8	16.6	17.2	20.6	22.6	22.8	46.73%
	Scen. III – Fixed W, changed B/S	6.2	8.2	12.4	16.8	16.2	18.0	25.6	37.08%
	Scen. III – Changed W	3.2	3.8	6.8	5.6	10.0	12.8	7.8	18.14%
	Scen. III – Total changes	9.4	12.0	19.2	22.4	26.2	30.8	33.4	55.22%
60%	Scen. II – Changed B and/or S	8.0	10.0	12.2	15.6	17.4	19.6	15.2	50.25%
	Scen. III – Fixed W, changed B/S	4.8	6.2	8.8	12.2	11.4	11.8	16.0	34.98%
	Scen. III – Changed W	3.0	3.0	5.8	6.0	9.4	11.4	6.6	21.83%
	Scen. III – Total changes	7.8	9.2	14.6	18.2	20.8	23.2	22.6	56.82%
40%	Scen. II – Changed B and/or S	5.8	7.0	9.4	12.2	10.6	13.4	11.6	54.12%
	Scen. III – Fixed W, changed B/S	3.0	3.6	6.8	7.4	7.2	7.6	11.2	34.16%
	Scen. III – Changed W	1.8	2.4	4.2	3.6	5.8	6.0	4.4	20.87%
	Scen. III – Total changes	4.8	6.0	11.0	11.0	13.0	13.6	15.6	55.03%

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# Conclusions

- Scalable column-generation based method for resilient VNet planning of time-varying traffic
- Our case study shows that:
  - Changing legacy traffic from one period to the next only really saves if we allow changing the working paths
  - ... but we need only to change around 20% of them
- Future work:
  - Optimization of choice of DC locations (e.g., ‘scattered’ vs ‘paired’, see ICTON 2013)
  - Extensive studies of different traffic patterns, over multiple periods

# Thank you ... any questions?

?

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